MODELING LOWER BOUNDS OF WORLD RECORD RUNNING TIMES

By

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ABSTRACT

MODELING LOWER BOUNDS OF WORLD RECORD RUNNING TIMES

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Fifty years ago, the first truly predictive model was developed for estimating a lower bound on a world record (the men's one mile record). This sparked an interest in modeling world record running times which has, in the past decade, faded. Then in 2008 at the Beijing Olympics 21 year old Usain Bolt broke the world record for both the 100 meter and 200 meter sprints. With his remarkable performance, Bolt not only established himself as one of track and field's premier athletes, but he reignited the debate over how low the world records for these events will ultimately fall. This paper presents several mathematical models for estimating a lower bound on world record running times for various distances ranging from the 100 meter sprint to the marathon. The models will use regression and order statistics to model the lower bounds.

For equations that require a solution to a system of equations that is not in closed form, we will be using numerical analysis methods on Mathematica® to estimate parameters. We will also use Mathematica® in solving linear systems for polynomial equations. These models will help us gain some insight into which records, if any, have already reached their pinnacle, and which records we can expect to be broken in the near future.

CHAPTER 1 INTRODUCTION

1.1 BEGINNING MODELS

In 1913, the International Amateur Athletic Foundation was founded as a committee that would keep official statistics on running records. However, even before the IAAF was founded, mathematicians had begun constructing models to predict how low certain running times would fall. In 1906 Kennely [7] used time-distance analysis to plot time based on a given distance for distances ranging from 100 yards to 10 miles. He let d be the distance of the race, in meters, and t the time the race took, in seconds. This yielded a log-log equation of:

$$\log(t) = 1.125\log(d) - 1.236$$

which, when solved for time *t*, gives:

$$t = 0.0581d^{1.125}$$

Though Kennelly did not construct this model strictly for predictive purposes, he did make a prediction for the ultimate world record in the mile run, which he estimated to be 3:58.1. The current world record for the men's one mile is 3:43.13.

1.2 PREDICTIVE MODELS

The first mathematical model constructed strictly for predictive purposes was done in 1958 by Lucy [8], who used the year to derive an exponential function for predicting the ultimate time in the one mile run, where t_n is the time, in seconds, for the mile record in the year (1950 + n), and *k* and *a* are estimated parameters:

$$t_n = t_0 + ka^n$$

Using this model, Lucy predicted an ultimate best time for the mile of 3:38, which is close to the record time now, fifty years later, of 3:43.13.

Just a couple of years after Lucy developed his model, Deakin [5] developed two distinct models using data from the years 1911 – 1965 in which world records were broken. For both models t_n is the time for a given year n years after 1911 in which a world record was set, α and β are parameters estimated through iterative methods, and γ (and in the second equation δ) are parameters estimated by looking at the plot and guessing the best fit curve:

$$t_{n} = \alpha e^{-\beta (n-1)} + \gamma$$
$$t_{n} = \alpha - (2 \beta/\pi) \arctan (\gamma (n-1) + \delta)$$

While it is clear that the arctan function was used because of its decaying shape, it seems illogical to use this equation because there is no reason to believe that running time data would behave according to an arctan function.

The second equation was also used for backward extrapolation to estimate an upper bound on what the fastest humans have always been capable of running in the mile, which Deakin estimated to be 5:00. Though both equations were crudely derived, both yielded the same lower estimate for the mile of 3:34, which is very close to Lucy's prediction of 3:38, and nine seconds from the current world record of 3:43.13.

In 1983 Schultz and McBryde [12] devised four unique models for predicting record performances in track and field. These models were derived by examining the best performance of the year for every year from 1900-1982. Subsets of these years were also fitted to the models to determine the effects of world events, such as WWI and WWII, but no significant influences were found. The authors found that all four models were approximately equally effective in fitting data from the shorter distances, but the model that best fit longer distances was an exponential model where t_n is the predicted time for a given year that is n years after 1899, α is the limit of t_n as $n \rightarrow \infty$, and β and γ are estimated parameters, with the restriction that γ must be negative:

$$t_n = \alpha + \beta e^{\gamma t}$$

The unique characteristic about this equation is that it models all track and field events, not just the running events. The only change to be made between running events and field events is the sign on β , which is positive for running events. Using their equation, Schutz and McBride estimated the ultimate mile time to be 3:28.5, compared to the current world record of 3:43.13 today.

In 1993 Schutz and Liu [13] reformatted this equation to fit a larger, more current sample, and used it for predictive purposes of one event, the men's 1500M. This new model, called the Random Sampling Model, works under the assumption that improvement has stabilized, and that a new record would be an outlier of a data sample. We will consider this idea in Chapter 3.

Chatterjee and Chatterjee [4] developed a model to predict ultimate times based on past Olympic performances. They used the winning times for the 100m, 200m, 400m and 800m races at the Olympics from 1900 - 1976 to look at the effects of Olympic year and race distance on time. In their equation, *i* refers to the *i*th Olympics within the range (1900 - 1976) and *j* refers to the *j*th event, with the events numbered 1 - 4 from shortest to longest distance. The variable *d* represents the actual distance of the race, measured in meters:

$$log(t_{ij}) = 0.249 + 0.150^{-0.069j} - 0.00263 \ d_j + 0.229 \sqrt{d_j}$$

Though the model was developed for the prediction of lowest times, the authors were more interested in predicting the times at the next Olympics, which would have been the 1984 Olympics in Los Angeles. The lowest times for the four events were not reported, but their predictions and the observed times for the 1984 Olympics were:

	0	
Race	Prediction	Actual
100m	9.98	9.99
200m	19.80	19.80
400m	44.77	44.27
800m	104.62	103.00

 Table 1: Chatterjee Estimates and Actual Times

Whereas Schutz and McBryde's model applied to all events from track and field, Morton [10] developed a model to predict times which included a unique factor in shorter track events: lane allocation. Time was determined by factoring the effects of gender (G), distance of the race (D), lane (L), and hurdle effects (H) (which is obviously not applicable in all races), and the equation Morton arrived at is:

 $T = 21.71 + 26.75D + 4.52H - 1.87G - 2.64DG - 1.48HG + 0.37DL - 0.30HL - 0.18GL + 0.84HL^2$

With this equation, Morton analyzed world record performances in reference to lane allocation. His study found no significant pattern in lane allocation and world record performance for the events, suggesting that for the races where lane assignment is mandatory throughout the entire race, lane allocation does not significantly affect performance. His results contradict the two theories that the outer lanes have an advantage due to the more gradual curvature, and the inner lanes have an advantage because of the ability to see the other competitors.

1.3 ALTERNATIVE MODELS

In addition to predictive mathematical models, where the main focus is to minimize time over a given distance, there have been several attempts to formulate a predictive model on the upper bound of maximum speed over a given distance.

One of the most prominent studies for this subject was done by Joseph Keller [6] in 1973. In trying to determine the optimal strategy for pacing in a race of a given distance so that time is minimized, Keller examined oxygen flow and the amount of energy provided from this flow, resistive forces per unit mass, force exerted by a runner, and velocity to calculate the maximum distance such that maximum velocity can be maintained. Keller calculated this value to be 291 meters. This suggests that for all distances in our study, with the exception of the 100M and 200M sprints, a pacing strategy must be utilized to minimize time.

As a result of this discovery, equations for pacing strategies have been developed to minimize the effects of fatigue during a race of a longer distance. Frank Mathis [9] examines Keller's optimal velocity equation and modifies it to include a variable for oxygen deficit during a race.

1.4 PREVIEW

The previously mentioned studies all contribute to the idea of finding a lower bound on running times. Some of the models extrapolate known data to predict future outcomes. In predicting, there has been a wide range of ultimate record times – some of which have been already broken, some of which will seemingly never be broken, and some of which are just small percentage points away from being broken. Other models explain the effect some variables have on running time or velocity, including resistive and propulsion

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forces [2], VO_2 max, fatigue [3,6], pacing [1], and oxygen consumption. Though each contribution is different, all of the models will be considered as a unique equation is constructed for predicting a lower bound on world record running times.

CHAPTER 2 UPDATING THE EQUATIONS

2.1 GATHERING THE DATA

We used the International Amateur Athletic Foundation to gather our data [16], which includes exclusively men's data as of now. We have compiled three classes of data: The first list is a list of the best recorded times in history for a given event. The next list is a list of the best time of the year for the event, and the final list is a list of the world record progression for the event. The world record list includes only the times that better the previous world record; the instances when the world record time was matched for an event were not included in the data set. The following figure gives the number of data points gathered for each event:

Table 2. Rumber of Data I offits							
Event	# Data Overall Best	# Data Best of Year	# Data WR Prog.				
100m	555	34	19				
200m	264	34	21				
400m	178	33	22				
800m	167	32	17				
1500m	253	29	33				
One Mile	178	32	31				
5,000m	296	22	35				
10,000m	160	25	38				
Half-Marathon	170	19	17				
Marathon	152	18	35				

 Table 2: Number of Data Points

It is to be noted that, for the best time of the year data, there are some years for which the best time was not available. For the world record times, some data sets have multiple values for the same year, due to the fact that the world record was broken more than one time in the same year. There are also many instances when there are several year gaps in the world record data. For example, the men's one mile world record was broken three times in 1981, but has only been broken three times since then, most recently in 1999.

As of now, we are only looking at the data from 5 events: 400m, 800m, One Mile, 5000m, and Marathon. The other distances will be examined at a later time.

2.2 EXAMINING LUCY'S EXPONENTIAL EQUATION

The first family of equations we are looking at is the exponential decay function. We focus on the form displayed by Lucy's equation. Recall that t_n is the best time of the year, and is known data; t_0 is the lower bound on the time; n is the number of years after 1950; and k and a are parameters to be estimated:

$$t_n = t_0 + k\alpha^n$$

Using the data from 1950 - 1958, Lucy found a lower bound for the men's one mile time of 218 seconds, or 3:38, which is approximately five seconds faster than the current world record of 223.13 seconds, or 3:43.13.

Using this general form of the equation, we use the best mile times from the years 1967 - 2007 to estimate the lower bound, t_0 . We have reindexed the value of n to be the number of years since 1967 for our version of the model. We also will estimate the parameters k and a.

To find these estimates, first we define an equation to find the sum of squares for the data points:

$$SS = \sum_{n=0}^{40} (t_n - t_0 - k\alpha^n)^2$$

We are trying to minimize the value of the sum of squares equation, so we take the partial derivative of the equation with respect to each variable and set each of these derivatives equal to zero. We end up with a system of three equations and three unknowns, so now we solve for the unknowns. Unfortunately, the system of equations does not have a closed-form solution. Thus, we must use a numerical method to approximate a solution. We use the FindRoot function on Mathematica to approximate the values for the parameters. Because this is a numerical process, we are required to initialize values for the parameters. We started by using the estimates Lucy was able to find in his equation:

$$SSA = \partial_{\alpha}SS \qquad SSK = \partial_{k}SS \qquad SST = \partial_{t}SS$$

FindRoot[{SSA = = 0, SSK = = 0, SST = = 0}, {a, 0.933}, {k, 31}, {t, 218}]

However, new estimates were not able to be found after 10,000 iterations, so the original estimates were adjusted. The value for k tended to be much lower than Lucy's estimate as it was adjusted, and the value for a tended to be slightly higher, so the original guesses were adjusted accordingly. Note that the value for a must be less than one in order for the equation to approach a lower bound. Using the new initial values, the numerical approximation converged in 9 iterations:

$$SSA = \partial_{\alpha}SS \qquad SSK = \partial_{k}SS \qquad SST = \partial_{t}SS$$

FindRoot[{SSA = = 0, SSK = = 0, SST = = 0}, {a, 0.94}, {k, 10}, {t, 218}]
{a = 0.947, k = 6.079, t = 225.89}

Our value for the lower bound, 225.89 seconds, is approximately eight seconds slower than the bound Lucy found. Our value for a is approximately the same, but our value for k is significantly different. This difference can be explained by the role the

parameter plays in the equation. The parameters k and a determine the rate of decay for the function. Because our value for a is similar to Lucy's, we only consider the differences in k. In Lucy's equation, the domain was from [0, 8], whereas in our equation, the domain goes from [0, 40]. This allows for a more gradual decay, provided that the absolute decays of the two equations are comparable. For Lucy's equation, exact data could not be found, but based on estimates using the world record progression, we approximate the time decay from 1950 – 1958 to be ten seconds. The decay from 1967 – 2007 for our equation was approximately eight seconds. Thus the decay is more gradual, and the curve is less steep compared to Lucy's as shown by the lower k estimate in our equation.

If we replace one of the variables for which we solved into the sum of squares equation, we can graph the function for the other two variables to show that the minimum exists at the convergent values. So we plug t_0 into the equation and graph it:

Figure 1: Showing the minimum of the sum of squares graphically



This shows that the minimum occurs at the estimated parameters. Using these parameter estimates, we can calculate the sum of squares and variance for our data.

$$SS = 88.066$$
 $\sigma^2 = 3.037$

If we look at smaller subsets of the data, we can see how the times are behaving over different intervals of time and use this information to improve our overall model. So we repeat the sum of squares process for smaller subsets of the data, as well as for the world record progression:

Subset of Years	k	а	t_0	SS	Variance
1967 – 2008	6.078	0.947	225.89*	88.066	3.037
1967 – 1987	11.941	0.977	220.00*	12.263	1.226
1979 – 1989	0.742	0.198	228.25	12.999	1.625
1979 - 2000	11.300*	0.977	220.66	56.256	2.961
WR progression	1.800	0.037	237.20	2417.440	86.337

 Table 3: One Mile: Current World Record is 223.13 seconds (3:43.13)

In some instances (indicated by a *), the system of equations was not converging after 40,000 iterations. However, one of the parameters would stay approximately the same as the initial guesses were altered; thus that approximate value was substituted into the original equation, and the system became a system of two equations and two unknowns, which converged.

The exponential model fits relatively well for the subsets of years, but not for the world record progression. This trend was evident in the four other distances for which this method was used:

Subset of Years	k	а	t_0	SS	Variance
1968 - 2008	3.479	0.993	41.05	4.983	0.166
1968 - 1988	-0.490	0.001	44.35	1.594	0.159
1988 - 2008	2.497	0.351	43.79	2.3936	0.141
WR progression	2.662	0.208	45.14	18.790	1.566

 Table 4: 400m: Current World Record is 43.18 seconds

Subset of Years	k	а	t_0	SS	Variance
1973 - 2008	1.073	0.784	102.75	13.919	0.480
1973 – 1988	1.249	0.825	102.58	4.639	0.464
1993 - 2008	1.270	0.272	102.61	5.457	0.455
WR progression	0.297	0.031	106.19	233.264	16.661

 Table 5: 800m: Current World Record is 101.11 seconds (1:41.11)

Table 6: 5000m: Current World Record is 757.35 seconds (12:37.35)

Subset of Years	k	а	t_0	SS	Variance
1981 - 2008	30.446	0.959	755.34	900.699	47.405
WR progression	0.063	0.037	808.48	39165.100	1223.910

 Table 7: Marathon: Current World Record is 7439 seconds (2:03:59)

Subset of Years	k	а	t_0	SS	Variance
1985 - 2008	80.545	0.840	7580.22*	58,362.500	3,890.830
WR progression	3490.510	0.979	7098.28	89,286,820.000	2,880,220.000

Note that in the 400m data, the 1968 - 1988 subset has a *k* value of -0.49 and an *a* value of 0.001. Because the value of *a* is small, the equation approaches the lower bound quickly, and the sign on *k* becomes insignificant. The world record data has the worst fit for this model in all of the distances examined thus far. This should be expected though, as the model was made for the best time of the year data, not world record progression.

In many instances, the lower bounds estimated are actually slower than the current world record times. This does not mean that they are inaccurate though, as the variance in the data could account for the differences in the times. However, we look at a different model to see if there is a better fit and to see if we can get different lower bounds. The next model we examine is the power model.

2.3 POWER MODEL

The family of equations we look at in the power model has t_n as the best time of the year, t_0 as the lower bound on the world record time, n as the number of years after a given year, k and a as parameters to be estimated, and has the form:

$$t_n = t_0 + \frac{a}{(n+1)} + \frac{k}{(n+1)^2}$$

We use the same data that was used in finding the exponential equations. When finding the parameters and lower bound, we use the same process as in the exponential equation. The only difference is that this family of equations is polynomial and has a closed form solution. We can use a linear solve method, rather than a numerical approximation method, to solve for the parameters. We use the NSolve function in Mathematica to find the values of the parameters and the lower bound.

 Table 8: 400m: Current World Record is 43.18 seconds

Subset of Years	k	а	t_0	SS	Variance
1968 - 2008	-6.185	6.312	43.72	4.782	0.159
WR progression	-145.996	151.660	42.14	2.240	0.187

Subset of Years	k	а	t_0	SS	Variance
1973 - 2008	-1.945	3.010	102.62	12.117	0.418
WR progression	-199.454	211.441	99.91	17.152	1.225

 Table 9: 800m: Current World Record is 101.11 seconds (1:41.11)

Table 10: O)ne Mile: Curre	nt World Record	d is 223.13 secor	nds (3:43.13)
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Subset of Years	k	а	t_0	SS	Variance
1967 – 2007	-34.641	39.817	225.91	87.477	3.017
WR progression	-17,815.200	1,868.500	208.55	100.566	3.592

Subset of Years	k	а	t_0	SS	Variance
1981 - 2008	-49.836	68.099	766.39	1108.200	58.326
WR progression	-66,277.000	7,164.780	709.16	3906.020	122.063

Table 11: 5000m: Current World Record is 757.35 seconds (12:37.35)

 Table 12: Marathon: Current World Record is 7439 seconds (2:03:59)

Subset of Years	k	а	t_0	SS	Variance
1985 - 2008	-489.252	584.898	7527.12	46,589.500	3,105.970
WR progression	-285,875.000	58,118.000	7131.64	1,763,830.000	56,897.600

The most notable difference in the power model estimates, compared to the exponential model estimates, is that the power model estimates for the world record progression fit the data significantly better. The variances for each distance are lower, and in some cases comparable to the variances found in the best of year data.

Specific to the power model equations, all of the values for the parameter k have negative values. This gives an upper bound on the times when n is positive. Unfortunately, this is not something that is necessarily desired. The general shape of the graph of the equation shows that as n is small, the power function is not a good fit for our purposes, as it increases to the upper bound, then decreases similarly to an exponential function. This shows that we need to look into higher order power functions to fit the data, or perhaps look at another method to generalize the data.

Figure 2: Power function for the marathon best of year data showing the general shape of the power function graph.



We can compare the two families of graphs to each other to see if they are giving similar lower bounds, and we can also look to see in general which graphs fit the data better by comparing the variances for each model:

These fet comparing commences nom the exponential and power models							
Years	Var, Exp	Var, Power	% Diff.	$t_0 \operatorname{Exp}$	t_{θ} Power	% Diff.	
400m							
1968 - 2008	0.166	0.159	-4.20%	41.05	43.72	6.11%	
WR Prog.	1.566	0.187	-738.67%	45.14	42.14	-7.12%	
800m							
1973 - 2008	0.480	0.418	-14.89%	102.75	102.62	-0.13%	
WR Prog.	16.662	1.225	-1259.92%	106.19	99.91	-6.29%	
One Mile							
1967 - 2007	3.037	3.017	-0.67%	225.89	225.91	0.01%	
WR Prog.	86.337	3.592	-2303.86%	237.2	208.55	-13.74%	
5000m							
1981 - 2008	47.405	58.326	18.72%	755.34	766.39	1.44%	
WR Prog.	1223.910	122.063	-902.69%	808.48	709.16	-14.01%	
Marathon							
1985 - 2008	3,890.830	3,105.970	-25.27%	7580.22	7527.12	-0.71%	
WR Prog.	2,880,220.000	56,897.600	-4962.11%	7098.28	7131.64	0.47%	

Table 13: Comparing estimates from the exponential and power models

For all of the models except the 5000m best of year times, the power models are giving a smaller variance, which implies that the power models fit slightly better than the exponential models. This is illustrated most clearly in the world record progression data, where the variance in the power function models are significantly smaller than the variance in the exponential models; in one instance nearly 5000% better. This does not imply, however, that the power function is a great fit to the data; it simply implies that it fits better than the exponential function.

We now test to see if the power function is a good fit. We can do this by looking at the residuals of the data, which are defined as the differences of the data points and the respective function values for given *n* values. If a model truly is a good representation of the data, then the residuals will be independent identically distributed random numbers from a normal distribution with a mean of 0 and a variance of σ^2 . This means that a graph displaying the residuals and the values for *n* should look like a random scatter plot. So we consider one example: the marathon best of year times. Graphing the residuals on the y-axis and the *n* values on the x-axis, the graph looks like:





The sum of all of the residuals is -0.003; however, by looking at the graph, the data does not appear to be completely randomly scattered. This shows that the power function (at least for the marathon) may not be a good fit.

CHAPTER 3 DISTRIBUTION OF THE MINIMA

3.1 GAMMA DISTRIBUTION

The next method we examine involves the distribution of the minima of a data set. Given a domain of some *n* years after a set year, we look at each year individually. For each year, a subset of all of the times recorded for the given distance is analyzed. The analyzing process includes: attempting to fit the data to a distribution, calculating the sample mean and variance/standard deviation of the distribution, and finally finding some number $x \ni P[X < x] = p$; essentially *x* is the $100(1-p)^{\text{th}}$ percentile of the data. Repeating this analysis for each year in the data set where data is accessible will yield a series of similar distributions of which the means, μ_n , and the values of x_n will asymptotically decrease to some value t_0 , our lower bound.

This method was used to look at the marathon for the year 2008. As of right now, the subset consists of the men's Boston and the New York City marathon results, which yields a data set of 38,320 times. In analyzing the data, it was originally assumed that the data would be normally distributed, and a goodness of fit test would be used to verify this. However, upon constructing a histogram of the data, the normal distribution seems to be a bad fit:

Figure 4: Boston Marathon Men 2008



Figure 5: NYC Marathon Men 2008







Based on the charts, a normal distribution was rejected and a gamma distribution was examined as a possibility. The gamma distribution has density:

$$f(x) = \frac{x^{\alpha - 1}e^{\left(\frac{-x}{\theta}\right)}}{\theta^{\alpha}\Gamma(\alpha)}$$

Where

$$\Gamma(\alpha) = \int_0^\infty t^{\alpha - 1} e^{-t} dt = (\alpha - 1)!$$

3.1.1 METHOD OF MOMENTS

Using our data set, we use the Method of Moments (MoM) to estimate the parameters α and θ . Define μ'_r to be the r^{th} moment about the origin of the random variable *x* coming from a gamma distribution. By definition:

$$\mu'_r = E(X^r) = \int_0^\infty x^r f(x) \, \mathrm{d}x = \int_0^\infty x^r \frac{x^{\alpha - 1} e^{\left(\frac{-x}{\theta}\right)}}{\theta^\alpha \Gamma(\alpha)} \, \mathrm{d}x$$

Simplifying this expression and making the u substitution $u = (x / \theta)$ we get:

$$\begin{split} \mu'_{r} &= \int_{0}^{\infty} x^{r} \frac{x^{\alpha-1} e^{(\frac{-x}{\theta})}}{\theta^{\alpha} \Gamma(\alpha)} dx = \frac{1}{\Gamma(\alpha)} \int_{0}^{\infty} (\frac{x}{\theta})^{\alpha} (x^{r-1}) (e^{\frac{-x}{\theta}}) dx \\ &= \frac{1}{\Gamma(\alpha)} \int_{0}^{\infty} u^{\alpha} (u^{r-1}) (\theta^{r-1}) (e^{-u}) \theta du = \frac{\theta^{r}}{\Gamma(\alpha)} \int_{0}^{\infty} u^{\alpha+r-1} (e^{-u}) du \\ &= \frac{\theta^{r}}{\Gamma(\alpha)} \Gamma(\alpha+r) \end{split}$$

This allows us to generate equations for μ'_1 and μ'_2 :

$$\mu'_{1} = \frac{\theta}{\Gamma(\alpha)} \Gamma(\alpha + 1) = \frac{\theta \alpha!}{(\alpha - 1)!} = \theta \alpha$$

$$\mu'_{2} = \frac{\theta^{2}}{\Gamma(\alpha)}\Gamma(\alpha+2) = \frac{\theta^{2}(\alpha+1)!}{(\alpha-1)!} = \theta^{2}(\alpha+1)\alpha$$

Because we are dealing with a sample of the entire data set, we will be looking at the sample moments, defined as:

$$m'_{k} = (\frac{1}{n}) \sum_{i=1}^{n} (x_{i}^{k}) = \mu'_{k}$$

So to find the estimates for the parameters α and θ , we solve:

$$m'_1 = \mu'_1 = \theta \alpha$$

$$m'_2 = \mu'_2 = \theta^2 (\alpha + 1) \alpha$$

Giving us the equations:

$$\hat{\alpha} = \frac{(m'_1)^2}{m'_2 - (m'_1)^2} = n\bar{x}^2 / \sum_{i=1}^n (x_i - \bar{x})^2$$
$$\hat{\theta} = \frac{m'_2 - (m'_1)^2}{m'_1} = (\frac{1}{n\bar{x}}) \sum_{i=1}^n (x_i - \bar{x})^2$$

3.1.2 GOODNESS OF FIT

We can use these parameter estimates in the gamma function to run a Goodness of Fit (GoF) test on our data to check and see if a gamma distribution is reasonable for this data set. To use the goodness of fit test, we first divide our data set into bins, specifying that each bin must contain at least two data points. If a bin has less than two data points, it is combined with another bin; this combination incurs a loss of one degree of freedom for our test. Once we sort the data into our bins, we compute a chi-square test statistic based upon the comparison of the expected number of data (e_i) and the actual number of data (f_i) in each of the m bins:

$$\chi^{2} = \sum_{i=1}^{m} \frac{(f_{i} - e_{i})^{2}}{e_{i}}$$

We then compare this test statistic to a known chi-square statistic and reject the null hypothesis if:

$$\chi^2 \ge \chi^2_{\alpha,m-3}$$

where two degrees of freedom are lost from the parameter estimates, and one is lost from using a sample.

Looking at the data from the 2008 NYC marathon, we find that:

$$\hat{\alpha} = 26.441$$
 and $\hat{\theta} = 581.739$

and, using twenty-five bins, we compute a test statistic of $\chi^2 = 1101.509$, which is greater than the chi-square statistic $\chi^2_{.95, 22} = 12.338$, so we reject the null hypothesis that the NYC marathon data comes from a gamma distribution.

Looking at the data from the 2008 Boston marathon, we find that:

$$\hat{\alpha} = 32.356$$
 and $\hat{\theta} = 416.148$

and, using sixteen bins, we compute a test statistic of $\chi^2 = 1281.549$, which is greater than the chi-square statistic $\chi^2_{.95, 13} = 5.892$, so we reject the null hypothesis that the Boston marathon data comes from a gamma distribution.

Now if we combine the data from both of the marathons, we find that:

$$\hat{\alpha} = 25.126$$
 and $\hat{\theta} = 586.098$

and, using twenty-five bins, we compute a test statistic of $\chi^2 = 1992.497$, which is greater than the chi-square statistic $\chi^2_{.95, 22} = 12.338$, so, when using the method of moments, we reject the null hypothesis that the composite NYC and Boston marathon data comes from a gamma distribution.



Figure 7: Combined data frequencies (light) and the gamma expected frequencies (dark)

3.1.3 MAXIMUM LIKELIHOOD ESTIMATORS

One way to improve on the chi-square statistic is to use the MLEs instead of the MoM. Generally, the MLEs will yield a better estimate for the values of the parameters, which may help lower the chi-square statistic for the GoF test.

When calculating the MLEs, we use the same data set we were using when we calculated the MoM. Recall this gave us 38,320 data points. To get the MLEs, we first have to get our likelihood function, which is the product of our multiple gamma function values over our 38320 data points:

$$L(x) = \prod_{i=1}^{n} \frac{x_i^{\alpha-1} * e^{(-x_i/\theta)}}{\theta^{\alpha} * \Gamma(\alpha)}$$

We then take the natural log of the likelihood function in order to simplify the partial derivatives we will be taking next:

$$\ln L(x) = (\alpha - 1)\left(\sum_{i=1}^{n} - \ln x_{i}\right) - \frac{1}{n\theta}\left(\sum_{i=1}^{n} - x_{i}\right) - n\alpha \ln\theta - n\ln\Gamma(\alpha)$$

We then take the partial derivatives with respect to α and θ and set them equal to zero:

$$\frac{\partial \ln L(x)}{\partial \alpha} = A(\alpha) = \left(\sum_{i=1}^{n} \ln x_{i}\right) - n\ln\theta - \frac{n\Gamma'(\alpha)}{\Gamma(\alpha)}$$
$$\frac{\partial \ln L(x)}{\partial \theta} = \frac{1}{n\theta^{2}}\left(\sum_{i=1}^{n} x_{i}\right) - \frac{n\alpha}{\theta}$$

These equations reduce once we substitute the constants in that we know: the sum of the natural logs of the data, the sum of the data, and the number of data points. Making the substitutions, we get:

A (
$$\alpha$$
) = 367, 054.97 - n ln (θ) - n $\psi(\alpha)$

$$B(\theta) = \frac{564,319,465}{n\theta^2} - \frac{n\alpha}{\theta}$$

We use the Solve function on Mathematica to get the parameter estimates for α and θ . We can compare these MLE estimates with the MoM estimates and the chi-square they give:

	mparing	om byu	
	α	<u>0</u>	<u>x2</u>
MLE	26.87	548.11	2535.68
МоМ	25.13	586.10	2027.20

Table 14: Comparing Chi-Square Statistics

Notice that the MoM estimates give a smaller chi-square statistic than do the MLE estimates. This is significant because it is generally acknowledged that the MLE are better estimates than the MoM. However, both of these chi-square statistics are suboptimal, so we need another approach.

3.1.4 EXPANDING THE SAMPLE

Prior to now, we were looking at just two marathons- NYC and Boston. They are large marathons, and they give us a large sample size of data. However, this large sample size comes from looking at all of the runners, whereas we are only concerned with the "elite" runners; those that create the front tail of the distribution. In a marathon, there are different levels of runners in terms of ability. There are runners who run the race simply for the accomplishment of running a marathon. These are the runners who make up the right tail of our distribution. Also, there are runners who are semi-competitve; they may have been training for this race for many months. They are faster than those in the previous group, and they will make up the middle portion of the distribution. Finally, there are those who run the marathon competitively. They are trying (and are capable of) winning the race, or placing in the top few runners at the finish- these are the elite runners. Our results thus far have shown that a Gamma distribution may not be an appropriate fit for the entire distribution. However, we now consider the fact that each level of runners may follow a unique distribution, which may, when combined with the other levels, form a unique, unified Gamma distribution. While we had previously been studying the entire curve, we now look at the distribution of just the elite runners. Consequently, we truncate our data set.

Because we are now looking at a smaller subset of data, we also decide that we can get a better representation of the marathon race as a whole if we add more marathons. We take the top 1000 times from each of the five marathons in the marathon majors series, which include the Boston, NYC, Chicago, London, and Berlin marathons. We also include the Beijing marathon in this data set (2008), as is customary to include in the majors results in the year that it applies. For this project, we will assume that all marathons, in terms of course layouts, weather conditions, and altitudes, are created equal. This can be said because extremities in course and/or weather conditions, which consistently deviate from the norm in some marathons, will not yield any times within our data set.

To run a GoF test, we first group the data into bins. However, in doing this we notice a trend at the slower end of the spectrum:



Figure 8: Data from the Marathon Major Series

We find out the problem is in the Chicago marathon data. The times near the end of this marathon tended to be slower, thus causing only data from this marathon to be in the last set of bins. So rather than take the top 1,000 times from each marathon, we here decide to establish our truncation criteria at a race completion time of three hours. Over the six marathons, this gives us 4,365 data points.

We redistribute these 4,365 data points into bins, using a bin width of 200. With this histogram, we now solve for the gamma distribution parameters α and θ that minimize the chi-square statistic from our goodness of fit test. To do this, we utilize a guess and check, or a "grid search" method. We first pick a value for α , then find the θ value for this α value that minimizes the chi-square.

$$G(x) = \frac{x^{\alpha - 1} * e^{x} \frac{-x}{\theta}}{\theta^{\alpha} * \Gamma(\alpha)} \qquad [n * \int_{7400 + 200 * (c-1)}^{7400 + 200 * (c)} G[x] \, \mathrm{d}x \, , \{c, 1, 13\}]$$

The number of runners, for this 2008 data, is 107,670; though we only have 4,365 data points in our set, we are estimating the parameters of an entire distribution based on the tail, so we must use the n value for the entire distribution. Once we have done this for the θ values, we move to another α value and repeat, until we find the smallest chi-square statistic. With this method, we obtain parameter values of $\alpha = 40.0$ and $\theta = 363.5$, which when substituted give a chi-square statistic of 79.

We can graph the observed and expected frequencies using these parameter estimates, where the histogram displays our observed data, and the line displays the expected data:



Figure 9: Histogram of the Truncated Data Set

From the graph, notice that most of the error in our GoF comes from the last two bins. Because we are only concerned with the elite times, we find it appropriate to further truncate the data and eliminate these two bins, reducing our data set from 4365 to 2708 data points.

With this truncated data set, we re-estimate the parameters. Rather than use a grid search method, we discover that we can minimize the chi-square using Mathematica. We set up chi-square function, Ki2008, of α and θ :

Gams = Table[
$$n * \int_{7400+200*(c-1)}^{7400+200*(c)} G(x) dx, \{c, 1, 13\}$$
]

Data2008 = {10, 10, 24, 37, 30, 45, 56, 94, 110, 156, 245, 302, 410}

 $Chis2008 = ((Data2008-Gams)^2) / Gams$

$$Ki2008(\alpha, \theta) = Total[Chis2008]$$

We take the partial derivatives with respect to α and θ :

$$da = \partial_{\alpha} Ki2008(\alpha, \theta)$$
 $dt = \partial_{\theta} Ki2008(\alpha, \theta)$

We set them equal to 0 and use the FindRoot operation in Mathematica to solve, using our grid search estimates as the initial conditions:

FindRoot[{da == 0, dt == 0}, {
$$\alpha$$
, 40}, { θ , 363}]

 $\{\alpha \rightarrow 39.54\}$ $\{\theta \rightarrow 368.51\}$

Comparing these to the estimates from our grid search method, we can see the estimates are close:

Table 15: Comparing Estimates							
	<u>α</u>	<u>θ</u>	<u>x2</u>				
Grid Search	40.00	363.50	35.20*				
Mathematica	39.54	368.51	35.06				

*This value was calculated by substituting the grid search parameter estimates into the data set after the final two bins were eliminated.

Though the improvement is not substantial, we now know that the chi-square is minimized.

By repeating this process over the past seven years (through 2001) we can check for trends in the parameter values, which may give insight into the distribution that best fits the data from a given year:

Year	<u>a</u>	<u>θ</u>
2008	39.54	368.51
2007	38.96	383.55
2006	37.77	389.58
2005	32.89	464.80
2004	30.21	517.15
2003	35.51	425.06
2002	38.10	381.64
2001	33.44	447.49

Table 16: Parameter Values by Year

Running a regression analysis on Excel on these parameters, we find that there is no significant trend in the values of α or θ . So we decide that we need to look at the distribution from another perspective, and we turn to order statistics.

CHAPTER 4 ORDER STATISTICS

4.1. DERIVING THE MINIMUM ORDER STATISTIC

We know the distribution for the minimum of a set of data is described by the equation:

$$Y_{\min}(x) = n(1 - F(x))^{n-1}f(x)$$

where f(x) is the density function of the data, F(x) is the integral of this function, and n is the number of data points. For this application, our density function is the gamma distribution function.

Before we look at the distribution of the minimum, we again decide that it is appropriate to truncate the data further. We decide to cutoff the data at 10,000 seconds, reducing our sample size from 2708 to 1529. Running the data through the minimum chisquare method we used with Mathematica, we get parameter values of $\alpha = 37.17$ and $\theta =$ 398.14, which give a chi-square statistic of 27.28. Using these parameter estimates, we compute the formula for the distribution of the minimum of this data, which gives the xvalue (time in seconds) with the highest probability of being the minimum time from the given data set. The equation for the distribution of the minimum order statistic is:

$$Y_{\min}(x) = (2.39 * 10^{-144}) x^{38.54} (e^{-.0027x}) (-6.22 * 10^{-15} + (2.65 * 10^{-46}) \Gamma(39.54, .0027x))^{1528}$$

Plotting this equation with time in seconds on the x-axis and the probability of the respective x-value on the y-axis, we see:

Figure 10: Graph of the Probability Density Function



Using the partial derivative of the function we find the time with maximum likelihood to be 8153.66 seconds. Unfortunately, this is not an optimal estimate, as the actual minimum for our data set is 7439 seconds. We can repeat the same process on the previous six years and we get the estimates below. Note that n_1 is the number of total competitors in all of the marathons, n_2 is the number of competitors who have met our qualification time, α and θ are our parameter estimates that minimize the chi-square statistic for the GoF test, and *min* is the value, in seconds, that has the maximum probability of being the minimum of our data set according to our order statistics distribution of the minimum function and graph:

 ruste 177 Derryudon of Minimum by Teur										
Year	<u>n</u> 1	<u>n</u> 2	α	θ	<u>x2</u>	Min	Actual Min			
2008	107670	1529	37.17	398.14	27.28	8153.66	7439.00			
2007	106169	1050	40.22	370.90	29.50	8605.24	7466.00			
2006	103517	1475	36.08	412.62	25.23	8136.84	7556.00			
2005	103509	1255	28.70	553.49	50.56	8069.73	7622.00			
2004	100286	1198	26.44	614.82	53.23	8019.10	7576.00			
2003	101472	1168	35.51	425.06	57.00	8315.98	7495.00			
2002	94542	1523	35.70	413.96	39.26	8033.49	7538.00			

Table 17: Derivation of Minimum by Year

The values for the chi-square are not extremely small, but reasonable enough to use the gamma distribution to estimate the minimum. However, the values we are getting for our minimum from the order statistics are poor. The estimates are several minutes slower than the actual minimum of the data set which we are using to estimate this given minimum. This leads us to believe that the data set which we are using may not be an adequate representative sample. We have been going on the assumption that this sample would be a good representation of the whole due to the fact that we are picking the major marathons, the sample size is substantially large (greater than 100,000), and many of the best marathon times are run in these five (in some years six) marathons. Also, this data was what was available to us at the time of computation. However, a new database of data was discovered, and a sample was found which we believe may be a more representative sample.

4.2. A NEW SAMPLE

The new sample we have consists of the top times of the year for all marathons. Rather than giving a fixed number of data for each year, there is a cutoff time of 8280 seconds for this data set, and however many times were ran at least in that fast that year were included in the set. This gave data sets between 402 - 958 competitors per year.

We repeated the same process as before, including the estimation of the calculation of the parameters which minimize the chi-square, the derivation of the function that describes the minimum of the order statistics, and the approximation of the minimum time with the greatest probability. We did have to change our n_1 value for these calculations, as more marathons were included in this data set. Initially, we made a rough approximation of n = 200,000 and performed the iterations on the data set for 2008. This gave us parameter values of $\alpha = 35.57$ and $\theta = 369.09$, with a chi-square statistic of

140.15. Using these parameters to graph the distribution of the minimum order statistic, we get a graph of:



Figure 11: Graph of the Probability Density Function

This graph gives a value for time with the highest probability of being the minimum of 7321.33. Note that this method gives a minimum value with highest probability that is actually faster than each of our data points. Unfortunately, our chi-square statistic is exceedingly too large. So we re-evaluated our estimate for n to see if maybe we were over or underestimating the actual value. We obtained a list of all of the marathons that hosted the runners who had a time in our new, faster, smaller data set, and estimated the total number of runners in those races. There turned out to be 147 different marathons throughout the world in which a runner raced a time that was competitive enough to make our list. Unfortunately, a large number of these marathons were smaller, foreign marathons, and thus it was impossible to obtain participation numbers for all of the races. So we had to use an estimation method to try and approximate the total number of runners.

To find an accurate estimation, we divided the marathons into two different strata: "large" and "small" marathons. Large marathons we defined to have over 10,000 participants, and small marathons we defined to have less than 10,000 participants. It is known that all of the marathons in the major series are large marathons, but there are an additional few that are not included in the major series but that can be classified as large marathons. We found that, for 2008, there were approximately 15 large marathons and 132 small marathons. This seems accurate, as historically the greater majority of marathons are smaller events. Most of the participation number data was available for the larger data, due to the fact that larger marathons get more publicity, and thus the information about the races are more widely distributed. However, only some of the smaller marathons we found an average for each strata and summed the products of the respective averages and number of races, which resulted in a data set of approximately 600,000. This compares to the amount of approximately 1.1 million we would have gotten by taking the average without stratifying the data *a priori*.

Now that we have better approximated the total number of runners, we recalculate the parameters and chi-square statistic, and plot the minimum order statistics function with the new value for n (600,000). The new values are $\alpha = 37.36$, $\theta = 373.72$, with a chi-square of 115.16. The graph of the minimum order statistic can be seen below:





The graph yields a minimum value with highest probability of 7910.68, which is higher than was the value which was calculated when n = 200,000. This is a counterintuitive result; one would think that as the number of runners increased, the value of the minimum order statistic, or the lowest time ran out of all the runners, would decrease directly with the increase of runners. We are finding that this is not the case.

4.3. EQUATING THE SUMS

Another thing that we find is that, unfortunately, our chi-square statistic is not sufficiently decreasing as we increase our value for n. This leads us to believe that the gamma function may not actually be a good fit for marathon runners. However, it is here that we realize there is a flaw in one of our underlying assumptions about the chi-square GoF test. When this test is utilized, a group of data is usually compared to an entire curve. However for our purposes, we are only concerned with the front tail of the curve. Because of this, we cannot assume that the sum of the expected values will equal the sum of the observed values. This is something we need to go back and check in our first group of data, consisting of the marathons from the major series, as it may be the source of the surprising results from the minimum order statistic.

4.3.1. CHECKING THE ORIGINAL DATA SET

Looking back, we find that the sums of the observed and expected values are nearly the same, differing by about 1%, so minor alterations are done to the parameters. The original table is listed, followed by the table when we adjust α and recalculate the chi-square statistic and minimum, followed by the table when we adjust θ and recalculate the chi-square and the minimum, and finally the table when we adjust the value of n and recalculate the chi-square and minimum values:

41

			0				
Year	<u>n</u> 1	<u>n</u> 2	α	<u>0</u>	<u>x2</u>	Min	Actual Min
2008	107670	1529	37.17	398.14	27.28	8153.66	7439.00
2007	106169	1050	40.22	370.90	29.50	8605.24	7466.00
2006	103517	1475	36.08	412.62	25.23	8136.84	7556.00
2005	103509	1255	28.70	553.49	50.56	8069.73	7622.00
2004	100286	1198	26.44	614.82	53.23	8019.10	7576.00
2003	101472	1168	35.51	425.06	57.00	8315.98	7495.00
2002	94542	1523	35.70	413.96	39.26	8033.49	7538.00

Table 18: Original Parameter Values

Table 19: Adjusting α

Year	<u>n</u> 1	<u>n</u> 2	α	<u>0</u>	<u>x2</u>	<u>Min</u>	Actual Min
2008	107670	1529	37.18	398.14	27.34	8156.59	7439.00
2007	106169	1050	40.25	370.90	29.66	8613.63	7466.00
2006	103517	1475	36.10	412.62	25.35	8142.88	7556.00
2005	103509	1255	28.74	553.49	51.12	8085.33	7622.00
2004	100286	1198	26.48	614.82	53.85	8036.16	7576.00
2003	101472	1168	35.53	425.06	57.21	8322.23	7495.00
2002	94542	1523	35.73	413.96	39.51	8042.56	7538.00

Table 20: Adjusting θ

Year	<u>n</u> 1	<u>n</u> 2	α	<u>0</u>	<u>x2</u>	Min	Actual Min
2008	107670	1529	37.17	398.30	27.36	8156.94	7439.00
2007	106169	1050	40.22	371.23	29.65	8612.89	7466.00
2006	103517	1475	36.08	412.58	25.32	8142.09	7556.00
2005	103509	1255	28.70	553.38	51.00	8083.72	7622.00
2004	100286	1198	26.44	614.73	53.76	8034.98	7576.00
2003	101472	1168	35.51	425.39	57.26	8322.43	7495.00
2002	94542	1523	35.70	413.92	39.46	8041.78	7538.00

Table 21: Adjusting n

<u>Year</u>	<u>n</u> 1	<u>n</u> ₂	<u>α</u>	<u>0</u>	<u>x2</u>	<u>Min</u>	Actual Min
2008	107080	1529	37.17	398.14	27.38	8153.66	7439.00
2007	104766	1050	40.22	370.90	29.70	8605.24	7466.00
2006	102725	1475	36.08	412.62	25.34	8136.84	7556.00
2005	101630	1255	28.70	553.49	51.07	8069.73	7622.00
2004	98293	1198	26.44	614.82	53.83	8019.10	7576.00
2003	101451	1168	35.51	425.06	57.28	8315.98	7495.00
2002	93420	1523	35.70	413.96	39.51	8033.49	7538.00

Notice that as the difference between the observed sum and expected sum was not substantially large, the modifications to the parameters did not result in large changes in the chi-square statistics and minimum values.

4.3.2. CHECKING THE NEW DATA SET

We repeat this process on our current data set consisting of the best times of the year. Looking at the values we used for our parameters and for n, we see that the sum of the observed and the sum of the expected differ enough that recalculations are necessary. To do this, we decide to alter the process by making *n* a parameter along with α and θ . We now set up a system of three equations and three unknowns: the three partial derivatives of our chi-square function, each with respect to α , θ , and n, respectively.

Unfortunately, using the FindRoot function on Mathematica with this additional variable yields some problems. We could not get Mathematica to successfully converge to three numbers because it was not finding a sufficient decrease in the merit function. However, by using two sets of initial conditions that were substantially different (approximately 10% different for each parameter), we could see that the function was stopping the iterative process at approximately the same parameter values. We thus deemed it acceptable to use these estimates for our parameters, which were $\alpha = 38.03$, $\theta = 363.75$, and n = 560,600.

We substituted these parameter values back into our equation to first check and see if the sum of the expected was equal to the sum of the observed. We expected it to be close, as the values calculated are supposed to be the values that minimize the chi-square, which will be minimized when the difference between the sum of the expected and the sum of the observed is equal to zero. Substituting back in the gamma equation, we find that the sum of the expected is slightly different than the sum of the observed, so we adjust our parameters accordingly so that the sums would be equivalent:

1 able 22. A	ajusung a) Equate	ule Sullis
	<u>n</u>	α	<u>θ</u>
<u>Original</u>	560600	38.03	363.75
Adjusted	560392	38.23	363.34

Table 22: Adjusting to Equate the Sums

When substituting the adjusted parameters, we get a chi-square statistic of 116.13, which unfortunately is not sufficiently less than the chi-square statistic we had previously calculated. We assume the same problems will arise for the other years, and therefore conclude that the gamma distribution is suboptimal for approximating marathon runners, and we move on to other distances.

CHAPTER 5 SHORTER DISTANCES

5.1. REVISITING THE GAMMA DISTRIBUTION

Moving on from the marathon distance, there are nine other distances which we are examining: the 100m, 200m, 400m, 800m, 1500m, one mile, 5000m, 10000m, and the half-marathon. The plan was to use the gamma distribution to approximate lower bound for the world record in these distances, following the same general model as what we did with the marathon. Unfortunately, the gamma distribution did not adequately fit the marathon runners, so we are not sure how well it will work for the other distances.

The first measure we took in checking the appropriateness of the gamma distribution to the other distances was a simplistic "glance" test to see if the data looked like it could be approximated by a gamma distribution. The data sets we are looking at for these distances consist of the top approximately 200 - 400 times of the year for the given distance. Illustrated below are the graphs for the marathon, then all of the other distances in ascending distance order for 2008:









Figure 16: 400m Histogram













Figure 20: 5000m Histogram











Unfortunately, by simply looking at these graphs of the bins for the different distances, none of them seem to approach the shape of a gamma distribution. So we must now search for another way to describe this data.

5.2. ORDERING THE DATA

The first thing we do to examine the data is plot all of the times in an order versus time graph, starting with the 100m times from 2008:



Figure 23: 100m Times from 2008 in Ascending Order

The most noticeable feature of this graph is the kink around the 270th term. This kink turns out to be consistently noticeable in the data for each year we are looking at, and it appears at the same time for each year: 10.16 seconds. In the graphs for the other

distances, we are seeing the same occurrence. What this leads us to believe is that there is a cutoff time at this kink where the data changes distribution. From this point, we decide to look at only the data prior to this kink, as it appears to follows the same distribution, and we are not concerned with the slower times from the data set.

5.2.1. KOLMOGOROV – SMIRNOV TEST

Eliminating all of the slower times for each year in the 100m and plotting the remaining times together, we get:



Figure 24: 100m Times for All Years in Ascending Order

Notice that in this graph the times for 2008 seem to be faster as a whole than the times for the other years, indicating that times are getting faster in comparison to the previous years. However, there also seems to be more data for 2008, which may account for the difference. We can check to see if the difference is due to an advancement in performance or a larger data set by using a Kolmogorov - Smirnov test. This test measures the distances between data points in distributions to test whether or not the distributions are significantly different from one another. Performing this test for the

100m, we find that there is no significant difference between the data sets in the years 2002 - 2007, but the data in the year 2008 is significantly different from the rest. This tells us that the visible difference is not simply a factor of the larger sample size, and that times are improving as a whole in 2008.

If we look at the plots of the data in each year for all of the distances, the only other distance where we can see a difference large enough to warrant a Kolmogorov – Smirnov test is in the half-marathon, as seen in the graph:



Figure 25: Half-Marathon Times for All Years in Ascending Order

Running a Kolmogorov - Smirnov test on this data shows that the times from 2008 and 2007 are not significantly different than each other, but together they are significantly different than the times from the other years, indicating a recent improvement in half-marathon times.

Also worth noting is that there is no kink in the distribution of the times for this distance. The only other distance with no prominent kink is the one mile, which also has the most disperse graph of the times over the years:



Figure 26: Mile Times for All Years in Ascending Order

5.3. VARIABLES IN DECREASING TIME

Based on the graphs of the data over the years, and the use of the Kolmogorov – Smirnov test, we can see that some distances are showing the potential to have the world records fall (100m, half-marathon) while other show that the records seem to be in less jeopardy of falling (mile, 800m). So what we need to do is figure out a method for tracking the data over these past years and predicting how low (if at all) the record for a given distance will fall. Because our data set is limited in how far back we can go, it is only logical to assume that our extrapolations cannot be made too far into the future. It would seem impractical to make a prediction about what the world record time will be in the 100m 150 years from now. So for now, we will be making predictions about the world record time approximately 10 - 15 years into the future.

5.3.1. IDENTIFYING THE VARIABLES

To measure the how low the records will fall for each distance, we first need to identify the factors that are driving the times down. Studying the distances over the past several years, we conclude that there are five main variables that interact to drive the fastest times down. The first variable is what we are terming the "competition ratio", which we define to be "the number of times ran in a given year that are measured to be in the top twenty-five times ever ran up to that point in time". We can see that as this ratio approaches one, then the likelihood that a record will be broken will increase due to the competition increasing. The next variable is the increase / decrease in the competition ratio from the previous year. This is important because an increase or decrease by some given amount does not have the same implications for all values of the competition ratio.

The next three variables will not necessarily apply to all distances. We will call these three variables "incentive", "phenom", and "switch". Incentive is the most complex of the variables due to the many layers it has. We use the term layers to describe the many subfactors contributing to the overall incentive effect. Incentive is an increasing function of societal preference. As an event becomes more popular, there is more incentive to excel in that particular event. As the societal preference increases, the monetary rewards will increase. Monetary rewards include prize money, appearance fees, and endorsement deals. The dichotomy of the incentive effect can be contrasted in the extremes (100m and marathon) versus the middle distance races (800m, 1500m, mile). Trends have shown that the middle distance races are much less popular than they once were in society. In contrast, societal preference for the extremes in the distances are steadily increasing. The monetary effects of this trend are evident. The top two or three middle distance runners will receive approximately \$50,000 - \$75,000 as an appearance fee to run in a highly promoted professional race [15]. For the marathon, the top ten or twelve marathoners in the world receive appearance fees of approximately \$200,000; Paula Radcliffe, the female world record holder in the marathon and still an active marathon competitor, commands an appearance fee of \$500,000 [11].

On the other extreme, Usain Bolt, world record holder in the men's 100m and 200m, has a multi-million dollar contract with Puma and other sponsors, and receives appearance fees of approximately \$250,000 per race [17]. It is no coincidence that the marathon and the 100m are the two events most recently to have the world record broken (both in 2008).

One of the other variables we are considering is what we are calling the "phenom effect". This term can alter by the year, and refers to the presence of a runner who is currently running in a given event with the capability to lower the world record in that event within the next approximately 2 - 3 years. We will define the capability to break a world record as "having ran a time that is within 1% of the world record time in that year". This will also account for the current world record holder still actively competing in an event (as is the case with Usain Bolt in the 100m and 200m).

The final variable we will consider is what we are calling the switch effect. This is similar to the phenom effect, but it takes into consideration certain athletes who would normally qualify as phenoms in a given event, but who are not currently competing in that event. This effect was most recently demonstrated with an athlete named Haile Gebrselassie, who in 1998 broke the world record for both the 5000m and the 10000m. Gebrselassie would later switch to the half marathon and then the marathon, both of which events he successfully lowered the world record, and is currently the world record holder for the marathon. Of all the variables, this is the most subjective. In order to qualify for a switch effect, a runner must have publicly declared that they will be making a switch in an event. Also, switches must happen within similar events. The many distances we are covering can be divided into three categories: sprints, middle distances, and long distances. The sprints are the 100m, 200m, and 400m; the middle distances are the 800m, the 1500m, the mile, and sometimes the 5000m; the long distances are sometimes the 5000m, the 10000m, the half marathon, and the marathon. The reason the 5000m is not strictly defined is because history has seen some athletes make the jump between the 1500m and the 5000m, but it is a rare occurrence for an athlete to excel in both events.

5.3.2. QUANTIFYING THE VARIABLES

Now that the variables are defined, we can quantify them. The values we have found are listed below, with time measured in seconds. The rows highlighted in yellow indicate that a world record was broken in that year; the rows in orange indicate a world record was broken twice that year:

<u>100m</u>						
year	time	comp ratio	comp improve	incentive	phenom	switch
2008	9.69	0.44	0.20	1	3	0
2007	9.74	0.24	-0.08	1	2	0
2006	9.77	0.32	0.20	1	2	0
2005	9.77	0.12	-0.04	1	1	0
2004	9.85	0.16	0.16	1	4	0
2003	9.93	0.00	0.00	1	0	0
2002	9.89	0.00	-0.12	1	0	0

Table 23: Variables for 100m

<u>200m</u>						
year	time	comp ratio	comp improve	incentive	phenom	switch
2008	19.30	0.16	-0.04	1	0	1
2007	19.62	0.20	0.00	1	0	0
2006	19.63	0.20	0.20	1	0	0
2005	19.89	0.00	-0.04	1	0	0
2004	19.79	0.04	0.04	1	0	0
2003	20.01	0.00	-0.04	1	0	0
2002	19.85	0.04		1	0	0

Table 24: Variables for 200m

Table 25: Variables for 400m

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<u>400m</u>						
year	time	comp ratio	comp improve	incentive	phenom	switch
2008	43.75	0.12	0.04	1	0	0
2007	43.45	0.08	0.00	1	1	0
2006	43.62	0.08	0.08	1	1	0
2005	43.93	0.00	0.00	1	0	0
2004	44.00	0.00	0.00	1	0	0
2003	44.57	0.00	0.00	1	0	0
2002	44.45	0.00		1	0	0

 Table 26: Variables for 800m

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<u>800m</u>						
year	time	comp ratio	comp improve	incentive	phenom	switch
2008	102.69	0.00	0.00	0	0	0
2007	103.74	0.00	0.00	0	0	0
2006	103.09	0.00	0.00	0	0	0
2005	103.70	0.00	0.00	0	0	0
2004	103.06	0.00	-0.04	0	0	0
2003	102.52	0.04	-0.04	0	0	0
2002	102.32	0.08		0	0	0

Table 27: Variables for 1500m

<u>1500m</u>						
year	time	comp ratio	comp improve	Incentive	phenom	switch
2008	211.57	0.00	0.00	0	0	0
2007	210.54	0.00	0.00	0	0	0
2006	209.02	0.00	0.00	0	0	0
2005	209.30	0.00	-0.08	0	0	0
2004	207.40	0.08	0.04	0	2	0
2003	208.40	0.04	-0.12	0	0	0
2002	206.89	0.16		0	2	0

mile						
year	time	comp ratio	comp improve	Incentive	phenom	switch
2008	229.38	0.00	-0.04	0	0	0
2007	226.91	0.04	0.04	0	0	0
2006	230.32	0.00	0.00	0	0	0
2005	227.97	0.00	0.00	0	0	0
2004	229.84	0.00	0.00	0	0	0
2003	228.17	0.00	0.00	0	0	0
2002	228.28	0.00		0	0	0

Table 28: Variables for One Mile

Table 29: Variables for 5000m

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<u>5000m</u>						
year	time	comp ratio	comp improve	Incentive	phenom	switch
2008	770.18	0.04	-0.08	0	0	0
2007	769.53	0.12	-0.04	0	0	0
2006	768.09	0.16	0.00	0	0	0
2005	760.18	0.16	0.04	0	1	0
2004	757.35	0.12	-0.16	0	2	0
2003	768.81	0.28	0.28	0	0	0
2002	775.85	0.00		0	0	0

Table 30: Variables for 10,000m

<u>10000m</u>						
year	time	comp ratio	comp improve	incentive	phenom	switch
2008	1585.97	0.04	-0.08	0	1	0
2007	1606.19	0.12	0.00	0	0	0
2006	1595.63	0.12	-0.08	0	0	0
2005	1577.53	0.20	0.08	0	2	0
2004	1580.31	0.12	-0.08	0	1	0
2003	1589.22	0.20	0.00	0	2	0
2002	1609.38	0.20		0	0	0

Table 31: Variables for Half-Marathon

<u>half</u>						
year	time	comp ratio	comp improve	Incentive	phenom	switch
2008	3555.00	0.16	-0.48	1	0	0
2007	3513.00	0.64	0.44	1	9	0
2006	3535.00	0.20	0.08	1	5	1
2005	3556.00	0.12	0.08	1	2	0
2004	3611.00	0.04	0.04	1	0	0
2003	3621.00	0.00	-0.16	1	0	0
2002	3585.00	0.16		1	2	0

marathon										
year	time	comp ratio	comp improve	Incentive	phenom	switch				
2008	7439.00	0.32	0.24	1	4	0				
2007	7466.00	0.08	0.00	1	1	1				
2006	7556.00	0.08	-0.04	1	1	0				
2005	7580.00	0.12	0.08	1	0	0				
2004	7574.00	0.04	0.04	1	0	0				
2003	7495.00	0.00	-0.16	1	9	0				
2002	7538.00	0.16		1	10	1				

 Table 32: Variables for Marathon

We can see that the world records have most recently been broken in the shorter and longer events, but not in the middle distance events. This is reflected in the low values for our variables in the middle distance events, and the high values in the shorter and longer events.

CHAPTER 6 FUTURE WORK

6.1. GENERALIZED GAMMA DISTRIBUTION

While we only studied the Gamma distribution using two parameters, there is a generalized gamma distribution that has three parameters. Though significantly more complicated than the two parameter version, the generalized gamma may be able to produce parameter values that will give a chi-square statistic low enough that an approximation of running times would be acceptable with this distribution.

6.2. SHORTER DISTANCES

Unfortunately, there was not enough time this year to fully examine the effect of the variables on the fastest yearly times. We have gathered all of the data, and are still in the process of analyzing the data for trends in the variables to see how accurately we can model the times. While it is evident that there is a strong correlation between high values for the variables and fast times for a given year, we need to model the changes in the variables so we can predict the lower bounds on the times for the years. The variable "switch" will have to be adjusted yearly, as it would be impossible to predict. However, estimates on times can be made that take into consideration all scenarios for this variable.

The biggest hurdle in modeling these variables will be in predicting the cycles for certain distances. The middle distances were once the most popular distances for racing, but now are the least. The one mile has an advantage in the incentive effect due to it's' capability of reacting fastest to the societal preference if the interest is ever reignited. It is a distance that is ingrained in us from middle school, as it is a commonly used standard distance for testing one's physical capabilities. Compare this to the 1500m, just 109 meters shorter in distance, and it is run in the Olympics (the one mile is not). However,

the 1500m would take much more time to build up societal preference. All of the distances would react in a unique way to societal changes, and this is another factor that will be modeled in the variable estimation in the future.

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BIOGRAPHICAL SKETCH



Kevin Heisler is a senior at Stetson University pursuing a major in Mathematics, with plans to pursue a career in Actuarial Science after graduation. He was born in Jupiter, Florida on October 31, 1986. His father, Michael Heisler, is a captain in Palm Beach County Fire Rescue, and his mother, Belinda Heisler, is a retired Registered Nurse. Kevin has an identical twin brother, Brian, who also attends Stetson University and is majoring in Education with a minor in Mathematics.

Kevin has lived in Jupiter his entire life and attended Jupiter Christian School from kindergarten through twelfth grade. While at Stetson, Kevin ran on the cross country team for three years. In his spare time Kevin enjoys playing poker, as well as many other various card games. A self-proclaimed skeptic, he enjoys reading about conspiracy theories and learning how things work and how they are made. Kevin enjoys studying numbers and once memorized the first 1,000 digits of pi in four days on a bet from his roommate.